

Network Security (NetSec)

IN2101 – WS 16/17

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Network InSecurity

Network “Security” offered by our Secret Services

Attacker Models

General Attacker Model

Attackers Limited by their Position in the Network

Security Goals

Security Goals Technically Defined

Threats

Threats Technically Defined

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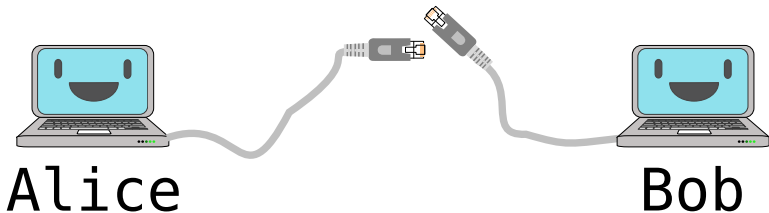
Threats

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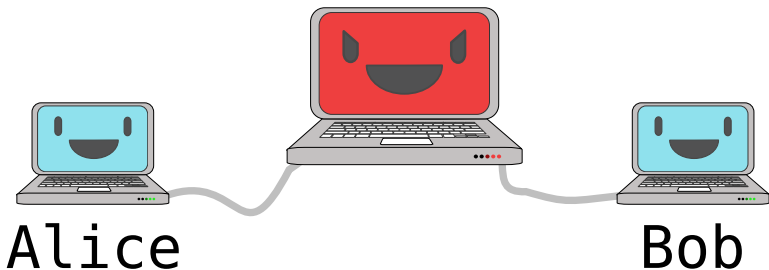
- By example: An Ethernet cable
- How secure is it?

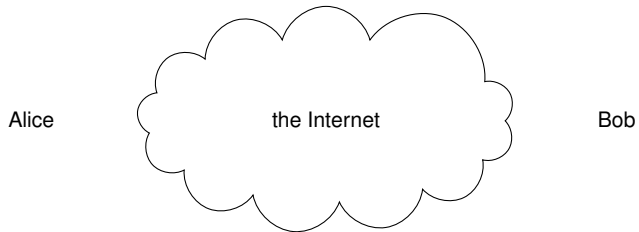


- Step 1: Obtain a knife
- Step 2: Add RJ45 adapters

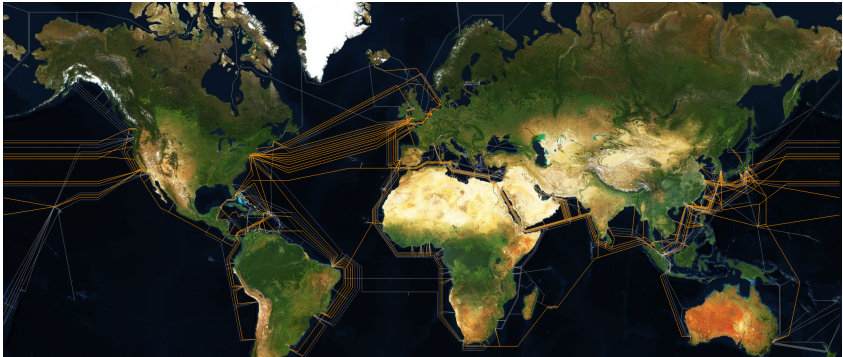


- Step 3: Configure transparent ethernet bridging
- You are now in full control of the traffic
 - read
 - modify
- Technical term: [Man in the Middle \(MitM\)](#)









<http://lifewinning.com/submarine-cable-taps/>

- Passive attacks: wiretapping, ...
- Active attacks: Quantum Insert, ...
- Combined: economic espionage, ...

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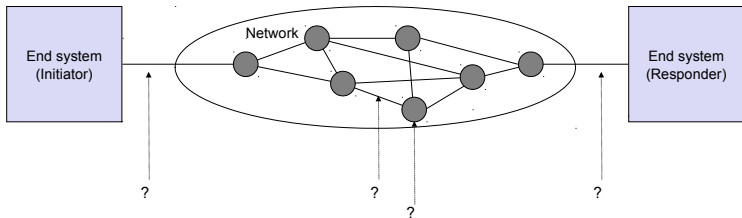
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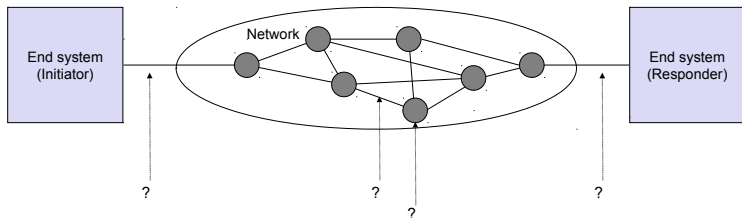
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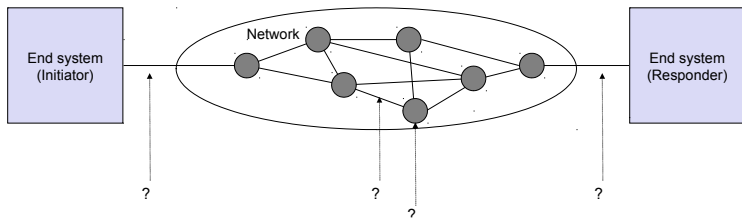
- Attacking communications on the message level
- Passive attacks:
 - Eavesdropping of messages
- Active attacks
 - all passive attacks
 - Delay
 - Replay
 - Deletion
 - Modification
 - Insertion

- The attacker **is** the network
- And can perform any active attack
- But cannot break cryptographic primitives
- This is called the **Dolev-Yao attacker model**
- If not stated otherwise, we will always assume this attacker model.

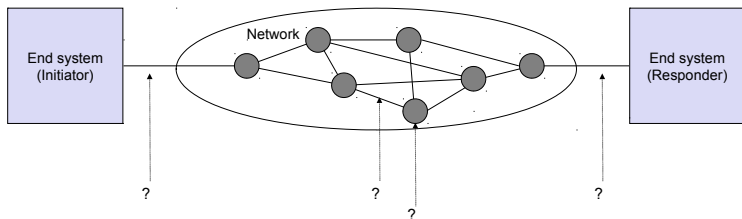




- Assume the Attacker is close to you
- Example: You sit in a cyber cafe and accidentally connected to the attacker's hotspot
 - The attacker can perform any active attacks on you
 - But you can bypass this attacker: Establish a **secure** tunnel to a server in the Internet
 - Route all your packets over the secure tunnel
 - The attacker can now perform only DOS (Denial Of Service) attacks against you



- Assume the Attacker is close to your servers
- Example: She rented a VM on the same host machine where your virtual server is running
 - The attacker could try to perform timing attacks against you
 - By measuring how long certain operations take at your server, the attacker might be able to break a security service
 - (only if the service is vulnerable to side channel attacks)
 - Such measurement is usually not possible over the Internet



- Assume the Attacker is somewhere in the Internet
- Internet: Best effort packet switching
- End-user has no control how packets are routed
- Are all AS/ISP trustworthy?
- Does your ISP alter your packets?
 - “value added service” i.e. your ISP places advertisement on the websites you are visiting
- NSA/GCHQ/BND/... black boxes are basically everywhere

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- Data Integrity
 - No improper or unauthorized change of data
- Confidentiality
 - Concealment of information
- Availability
 - Services should be available and function correctly
- Authenticity
 - Entity is who she claims to be
- Accountability german: „Zurechenbarkeit“
 - Identify the entity responsible for any communication event
- Controlled Access
 - Only authorized entities can access certain services or information

- What is needed to support non-repudiation? („*Nicht-Abstreitbarkeit*“)

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- What is necessary to support accountability?

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 - Authenticity

- What do you want to support deterrence („*Abschreckung*“)

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 - Accountability

- What is data origin integrity?

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 - Authenticity

- What is the difference?
- Authentication

- Authorization

- What is the difference?
- Authentication
 - Proves who you are
 - Associated security goal: Authenticity
- Authorization
 - Defines what you are allowed to do
 - Associated security goal: Controlled Access

- What is the difference?
- Authentication
 - Proves who you are
 - Associated security goal: Authenticity
 - E.g. your passport
- Authorization
 - Defines what you are allowed to do
 - Associated security goal: Controlled Access
 - E.g. “*are you on the VIP list?*”



My best attempt was registering to Black Hat with first name: "Staff" and last name: "Access All Areas"

<https://twitter.com/mikko/status/587973545797492738>

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- Abstract Definition
 - A threat in a communication network is any possible event or sequence of actions that might lead to a violation of one or more security goals
 - The actual realization of a threat is called an attack

- Masquerade
 - An entity claims to be another entity (also called “impersonation”)
- Eavesdropping
 - An entity reads information it is not intended to read
- Loss or Modification of (transmitted) Information
 - Data is being altered or destroyed
- Denial of Communication Acts (Repudiation)
 - An entity falsely denies its participation in a communication act
- Forgery of Information
 - An entity creates new information in the name of another entity
- Sabotage/Denial of Service
 - Any action that aims to reduce the availability and / or correct functioning of services or systems
- Authorization Violation:
 - An entity uses a service or resources it is not intended to use

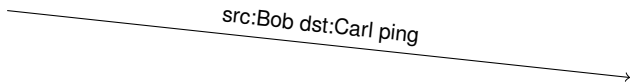
- Eavesdropping + Authorization Violation
- Example
 - Alice@Box\$./rootremoteshell \$ROUTER
root@router# tcpdump | grep password
- If Alice does not start modifying the traffic, she is a passive attacker
- Note: If not stated otherwise, we assume that attackers don't have remote code execution on our boxes

- Masquerade + Forgery of Information
- Example
 - Alice pretends to be Bob
 - Alice@Box\$ hping3 --count 1 --spooof \$BOB --icmp --icmptype 8 \$CARL
 - Bob gets an ICMP Echo Reply which he never requested
- Alice is an active attacker

Alice

Bob

Carl



- Alice: 192.168.1.170
- Bob 192.168.1.227
- Carl: 192.168.1.1
- Alice sends the spoofed packet
 - Internet Protocol Version 4, Src: 192.168.1.227, Dst: 192.168.1.1; ICMP Echo Request
- Carl replies to the source address specified
- Bob receives a lonely echo reply
 - Internet Protocol Version 4, Src: 192.168.1.1, Dst: 192.168.1.227; ICMP Echo Reply

```
192.168.1.1      192.168.1.227      ICMP      60 Echo (ping) reply      id=0xce1f, seq=0/0, ttl=61
```

- Denial of Service
- Example
 - Bob runs a webserver (http, tcp port 80) with very few memory
 - Alice floods Bob with TCP SYN packets
 - Alice@Box\$ hping3 --fast --count 42 --syn --destport 80 \$BOB
 - Bob allocates memory to store the 42 connections in the SYN-RECEIVED state
- Now Alice starts to deny that she is responsible for the attack
- Denial of Service + Forgery of Information + Denial of Communication Acts
- Example
 - Alice@Box\$ hping3 --fast --count 42 --rand-source --syn --destport 80 \$BOB
 - --rand-source: random spoofed source IP address

Capturing from Ethernet [Wireshark 1.12.4 (v1.12.4-0-gb4861da from master-1.12)]

File Edit View Go Capture Analyze Statistics Telephony Tools Internals Help

Filter: Expression... Clear Apply Save

No.	Time	Source	Destination	Protocol	Length	Info
736	686.042764000	56.10.51.117	192.168.1.227	TCP	60	1350-80 [SYN] Seq=0 win=512 Len=0
737	686.129344000	38.36.23.85	192.168.1.227	TCP	60	1351-80 [SYN] Seq=0 win=512 Len=0
738	686.229507000	36.116.117.78	192.168.1.227	TCP	60	1352-80 [SYN] Seq=0 win=512 Len=0
739	686.329714000	189.139.51.172	192.168.1.227	TCP	60	1353-80 [SYN] Seq=0 win=512 Len=0
740	686.429848000	242.114.151.137	192.168.1.227	TCP	60	1354-80 [SYN] Seq=0 win=512 Len=0
741	686.530802000	255.124.118.119	192.168.1.227	TCP	60	1355-80 [SYN] Seq=0 win=512 Len=0
742	686.630208000	161.10.181.62	192.168.1.227	TCP	60	1356-80 [SYN] Seq=0 win=512 Len=0
743	686.730401000	9.205.193.205	192.168.1.227	TCP	60	1357-80 [SYN] Seq=0 win=512 Len=0
744	686.830479000	205.95.119.125	192.168.1.227	TCP	60	1358-80 [SYN] Seq=0 win=512 Len=0
745	686.930632000	238.97.119.210	192.168.1.227	TCP	60	1359-80 [SYN] Seq=0 win=512 Len=0
746	687.030809000	194.238.30.56	192.168.1.227	TCP	60	1360-80 [SYN] Seq=0 win=512 Len=0
747	687.130950000	111.148.162.200	192.168.1.227	TCP	60	1361-80 [SYN] Seq=0 win=512 Len=0
748	687.230995000	225.60.95.186	192.168.1.227	TCP	60	1362-80 [SYN] Seq=0 win=512 Len=0
749	687.331114000	124.161.110.246	192.168.1.227	TCP	60	1363-80 [SYN] Seq=0 win=512 Len=0
750	687.431808000	193.202.206.237	192.168.1.227	TCP	60	1364-80 [SYN] Seq=0 win=512 Len=0

- Why does the attack succeed?
- This is a good opportunity to refresh your knowledge about the TCP 3-way handshake

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